### **Discrete Mathematics**

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### Review

- Truth assignment
- Truth valuation
- Tautology
- Consequence

### Outline

• Tableau proof system

# Terminologies

- signed proposition
- entries of the tableau
- atomic tableau

# Tableau

### Definition (Tableaux)

A *finite tableau* is a binary tree, labeled with signed propositions called entries, such that:

- All atomic tableaux are finite tableaux.
- If τ is a finite tableau, P a path on τ, E an entry of τ occurring on P and τ' is obtained from τ by adjoining the unique atomic tableau with root entry E to τ at the end of the path P, then τ' is also a finite tableau.

If  $\tau_0, \tau_1, \ldots, \tau_n, \ldots$  is a (finite or infinite) sequence of the finite tableaux such that, for each  $n \ge 0, \tau_{n+1}$  is constructed from  $\tau_n$  by an application of (2), then  $\tau = \cup \tau_n$  is a *tableau*.

### Tableau

### Example

A tableau with the signed proposition  $F(((\alpha \rightarrow \beta) \lor (\gamma \lor \delta)) \land (\alpha \lor \beta)).$ 

## Tableau

#### Definition

Let  $\tau$  be a tableau, P a path on  $\tau$  and E an entry occurring on P.

- *E* has been *reduced* on *P* if all the entries on one path through the atomic tableau with root *E* occur on *P*.
- P is contradictory if, for some proposition α, Tα and Fα are both entries on P. P is finished if it is contradictory or every entry on P is reduced on P.
- **③**  $\tau$  is *finished* if every path through  $\tau$  is finished.
- $\tau$  is contradictory if every path through  $\tau$  is contradictory.

# Proof

### Definition

- A tableau proof of a proposition α is a contradictory tableau with root entry Fα. A proposition is tableau provable, written ⊢ α, if it has a tableau proof.
- Solution A tableau refutation for a proposition  $\alpha$  is a contradictory tableau starting with  $T\alpha$ . A proposition is tableau refutable if it has a tableau refutation.

### Definition (Complete systematic tableaux)

Let R be a signed proposition. We define the *complete systematic* tableau(CST) with root entry R by induction.

- Let  $\tau_0$  be the unique atomic tableau with R at its root.
- 2 Assume that  $\tau_m$  has been defined. Let *n* be the smallest level of  $\tau_m$  and let *E* be the leftmost such entry of level *n*.
- S Let τ<sub>m+1</sub> be the tableau gotten by adjoining the unique atomic tableau with root E to the end of every noncontradictory path of τ<sub>m</sub> on which E is unreduced.

The union of the sequence  $\tau_m$  is our desired complete systematic tableau.

#### Theorem

Every CST is finished.

### Proof.

Reduce the *E* level by level and there is no *E* unreduced for any fixed level.  $\Box$ 

#### Theorem

If  $\tau = \bigcup \tau_n$  is a contradictory tableau, then for some  $m, \tau_m$  is a finite contradictory tableau. Thus, in particular, if a CST is a proof, it is a finite tableau.

Proof.

By König lemma.

#### Definition

We define  $d(\alpha)$ , the degree of a proposition  $\alpha$  by induction.

- if  $\alpha$  is a propositional letter, then  $d(\alpha) = 0$ .
- 2) if  $\alpha$  is  $\neg \beta$ , then  $d(\alpha) = d(\beta) + 1$ .

The degree of a signed proposition  $T\alpha$  or  $F\alpha$  is the degree of  $\alpha$ . If P is a path in a tableau  $\tau$ , then d(P) the degree of P is the sum of the degree of the signed propositions on P that are not reduced on P.

#### Theorem

Every CST is finite.

#### Proof.

Every path is finite with  $d(P_{m+1}) < d(P_m)$ .

## Next Class

- Soundness theorem
- Completeness theorem