

# Discrete Mathematics

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# Review of Lattice

- Ideal
- Special Lattice
- Boolean Algebra

# Examples of Proof

- Zeno's paradox
- Zhuang Zi's paradox
- Gong Sunlong's "a white horse is not a horse"
- ...

How can you **persuade** yourself and the others?

# Examples of Proof

## Axiom

*The axiom of group theory can be formulated as follows:*

- (G1) For all  $x, y, z$ :  $(x \circ y) \circ z = x \circ (y \circ z)$ .*
- (G2) For all  $x$ :  $x \circ e = x$ .*
- (G3) For every  $x$  there is a  $y$  such that  $x \circ y = e$ . (right inverse)*

## Theorem

*For every  $x$  there is a  $y$  such that  $y \circ x = e$ . (left inverse)*

# What is Logic

- Premise
- Argument
- Conclusion
- Follow
- Proof

# History of Mathematical Logic

- Aristotle(384-322 B.C.): theory of syllogistic
- De Morgan(1806-71), Boole(1815-64), Schröder(1841-1902)
- Frege(1848-1925), Russell(1872-1970)
- Post(1897-1954), Gödel (1906-78), Henkin(1921-2006), Herbrand(1908-31)
- Robbinson(1930-); Beth and Smullyan
- Leibniz(1646-1716) and Hilbert(1862-1943)

# Introduction to Mathematical Logic

- First order logic
  - Propositional Logic
  - Predicate Logic
- High order logic
- Other type of logic
  - Modal logic
  - Intuitionistic logic
  - Temporal logic

# Introduction to Mathematical Logic

- Proof system
  - Axiom
  - Tableaux
  - Resolution
- Two Components
  - Syntax
  - Semantics
- Algorithmic approach



# Order

## Definition (Partial order)

A *partial order* is a set  $S$  with a binary relation  $<$  on  $S$ , which is *transitive* and *irreflexive*.

## Definition (Linear order)

A partial order  $<$  is a *linear order*, if it satisfies the *trichotomy law*:  $x < y$  or  $x = y$  or  $y < x$ .

## Definition (Well ordering)

A linear order is *well ordered* if every nonempty set  $A$  of  $S$  has a least element.

# Countable and Infinite

## Definition (Countable)

A set  $A$  is *countable* if there is a one-to-one mapping from  $A$  to  $\mathcal{N}$ .

## Definition (Finite)

A set  $A$  is *finite* if there is a one-to-one mapping from  $A$  to  $\{0, 1, \dots, n - 1\}$  for some  $n \in \mathcal{N}$ .

## Definition

- 1 If  $A$  is not countable, it is *uncountable*.
- 2 If  $A$  is not finite, it is *infinite*.

# Countable and Infinite

## Theorem

*Let  $A$  be a countable set. The set of all finite sequence of elements in  $A$  is also countable.*

## Proof.

We can formalize it as

$$S = \cup_{n \in \mathcal{N}} A^n = A^1 \cup A^2 \cup \dots \cup A^n \cup \dots$$

Construct a mapping from  $A^n$  to  $\mathcal{N}$ . □

## Definition (Tree)

A *tree* is a set  $T$  (whose elements are called nodes) partially ordered by  $<_T$ , with a unique least element called the *root*, in which the predecessors of every node are well ordered by  $<_T$ .

## Definition (Path)

A *path* on a tree  $T$  is a maximal linearly ordered subset of  $T$ .

## Definition (Properties of tree)

- 1 The *levels* of a tree  $T$  are defined by induction.
- 2 The  $0^{th}$  level of  $T$  consists precisely of the root of  $T$ .
- 3 The  $k + 1^{th}$  level of  $T$  consists of the immediate successors of the nodes on the  $k^{th}$  level of  $T$ .

## Definition (Properties of tree)

- 1 The *depth* of a tree  $T$  is the maximum  $n$  such that there is a node of level  $n$  in  $T$ .
- 2 If there are nodes of the level  $n$  for every natural number  $n$ , we say the depth of  $T$  is infinite of  $\omega$ .

## Definition (Properties of tree)

- 1 If each node has at most  $n$  immediate successors, the tree is  $n$ -ary or  $n$ -branching.
- 2 If each node has finitely many immediate successors, we say that the tree is *finitely branching*.
- 3 A node with no successors is called a *leaf* or a *terminal node*.

## Theorem (König's lemma)

*If a finitely branching tree  $T$  is infinite, it has an infinite path.*

## Proof.

- 1 If there is no infinite path, the tree would be finite.
- 2 Split the successors of the node into two parts. One with infinite successors and the other with finite successors.





## Definition

A *labeled tree*  $T$  is a tree  $T$  with a function (the labeling function) that associates some objects with every node. This object is called the *label* of the node.

## Definition (Segment)

- 1  $\sigma$  is an *initial segment* of  $\tau$  if  $\sigma \subset \tau$  or  $\sigma = \tau$ .
- 2  $\sigma$  is an *proper initial segment* of  $\tau$  if  $\sigma \subset \tau$ .

## Definition (Lexicographic ordering)

For two sequences  $\sigma$  and  $\tau$  we say that  $\sigma <_L \tau$  if  $\sigma \subset \tau$  or if  $\sigma(n)$ , the  $n^{\text{th}}$  entry in  $\sigma$ , is less than  $\tau(n)$  where  $n$  is the first entry at which the sequences differ.

One way to define a linear order based on given tree.

## Definition (left to right ordering)

Given two nodes  $x$  and  $y$ ,

- 1 If  $x <_T y$ , we say that  $x <_L y$ .
- 2 If  $x$  and  $y$  are incomparable in the tree ordering, find the largest predecessors of  $x$  and  $y$ , say  $x'$  and  $y'$ 
  - 1 If  $x'$  equals  $y'$ ,  $x \leq y$  if and only if  $x$  is left to  $y$  relative to  $x'$ .
  - 2 Otherwise  $x <_L y$  if and only if  $x' <_L y'$ .

# Next Class

- Language of proposition logic
- Formation tree
- Truth table
- Connectives